



CPS3232

Applied Cryptography

Lecture 1: Introduction and Course Overview

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About Me

- Course responsible and lecturer
- Ph.D. and postdoc in Computer Science, and MSc in Information Security
- 15+ years of software industry experience (cyber security, development, consultancy, etc.)

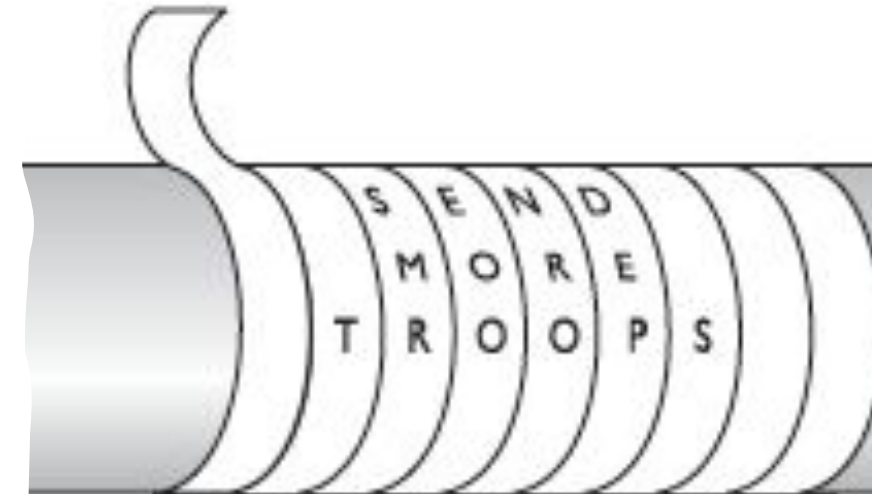
Learning Objectives

- Define cryptology and its types
- Explain the security services provided by cryptography
- Describe and analyze a few historical ciphers
- Identify and describe the differences between different types of attacks against ciphers
- Explain the general objectives and structure of the course, and expectations



Historical Ciphers

- Early signs of encryption in Egypt in ca. 2000 B.C.
- Historical ciphers include the Caesar cipher, Vigenère cipher, and Playfair cipher.
- Modern cryptographic systems are much more secure than historical ciphers.





What is Cryptography?

- *Cryptography* is the practice and study of techniques for secure communication in the presence of adversaries.
- *Applied cryptography* focuses on a subset of cryptographic constructions already practical and integrated into larger systems.

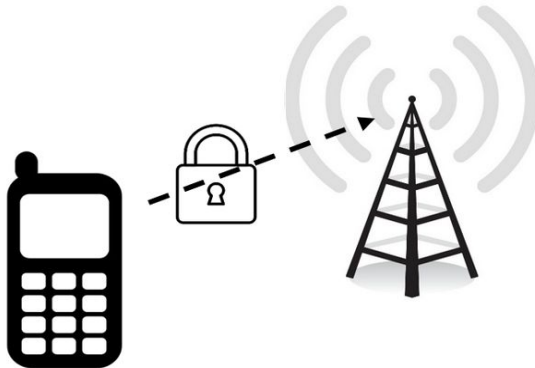
Applications of Cryptography



The Enigma Machine being used to secure radio communications during WWII

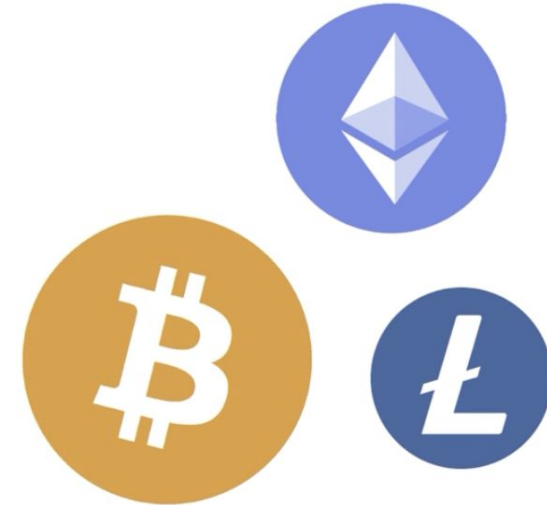


Secure online payments



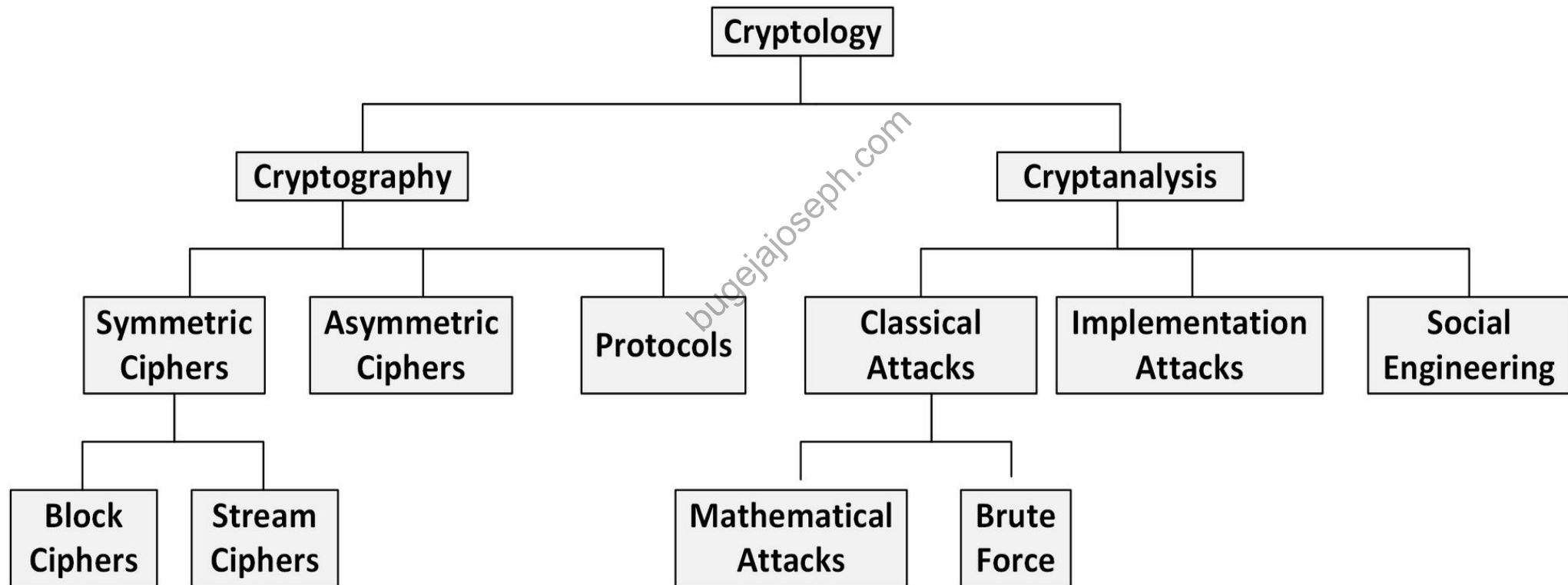
Social media posts

Can you think of other everyday applications of cryptography?



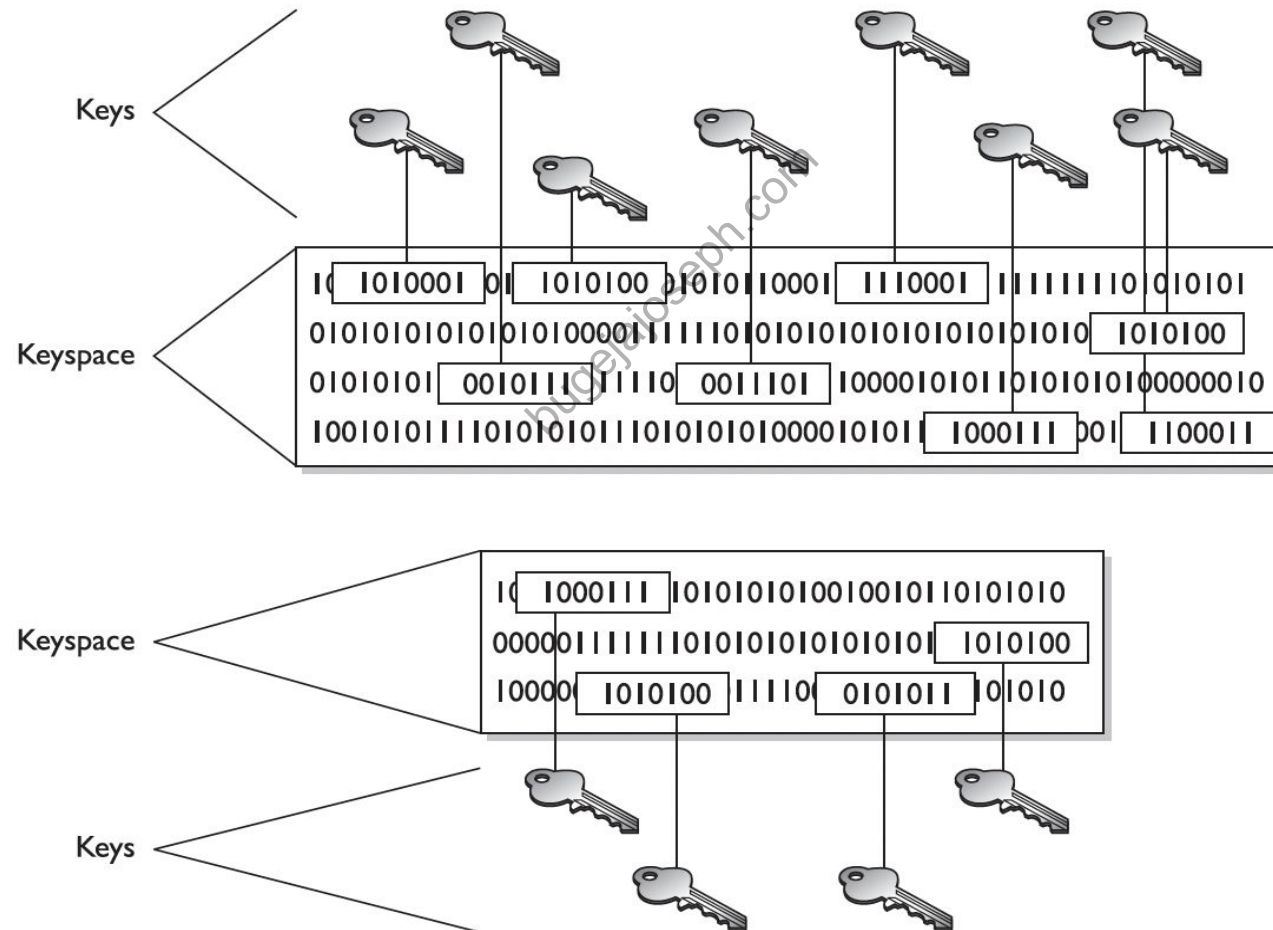
Cryptocurrencies

Overview of the Field of Cryptology



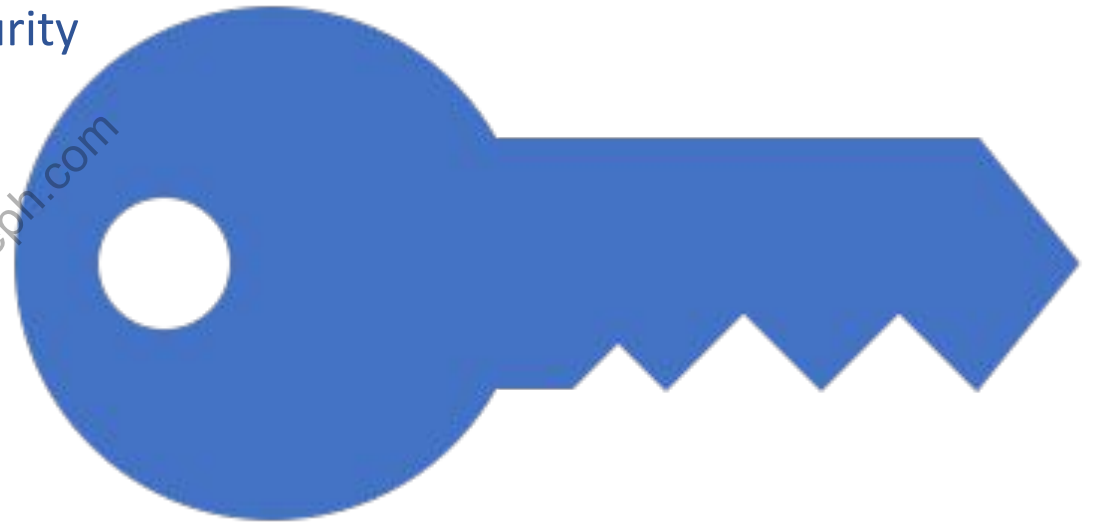
Keys and Keyspace

- A large keyspace allows for more possible keys, e.g., a key size of 512 bits would provide a keyspace of 2^{512} .



Kerckhoff's Principle

- There is in general no mathematical proof of security for any practical cipher.
- The only way to have assurance that a cipher is secure is to try to break it (legally) (and fail)!
- Kerckhoff's Principle is paramount in modern cryptography:
 - *A cryptosystem should be secure even if the attacker (Oscar) knows all details about the system, with the exception of the secret key.*
- To achieve Kerckhoff's Principle in practice, we should only use widely-known ciphers that have been cryptanalyzed for several years by good cryptographers!





Services of Cryptosystems

- **Confidentiality:** Renders the information unintelligible except by authorized entities.
- **Integrity:** Data has not been altered in an unauthorized manner since it was created, transmitted, or stored.
- **Authentication:** Verifies the identity of the user or system that created the information.
- **Non-repudiation:** Ensures that the sender cannot deny sending the message.

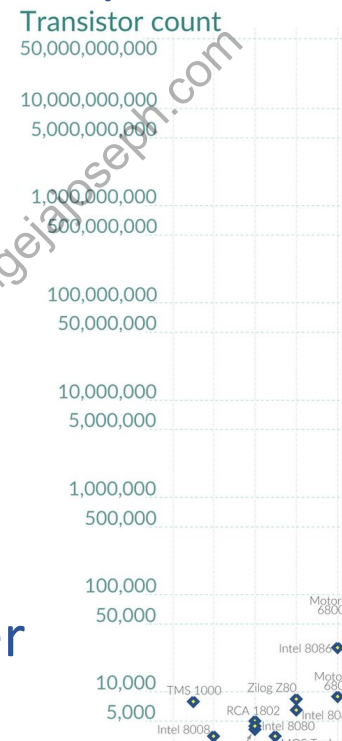
Can you think of another security service?

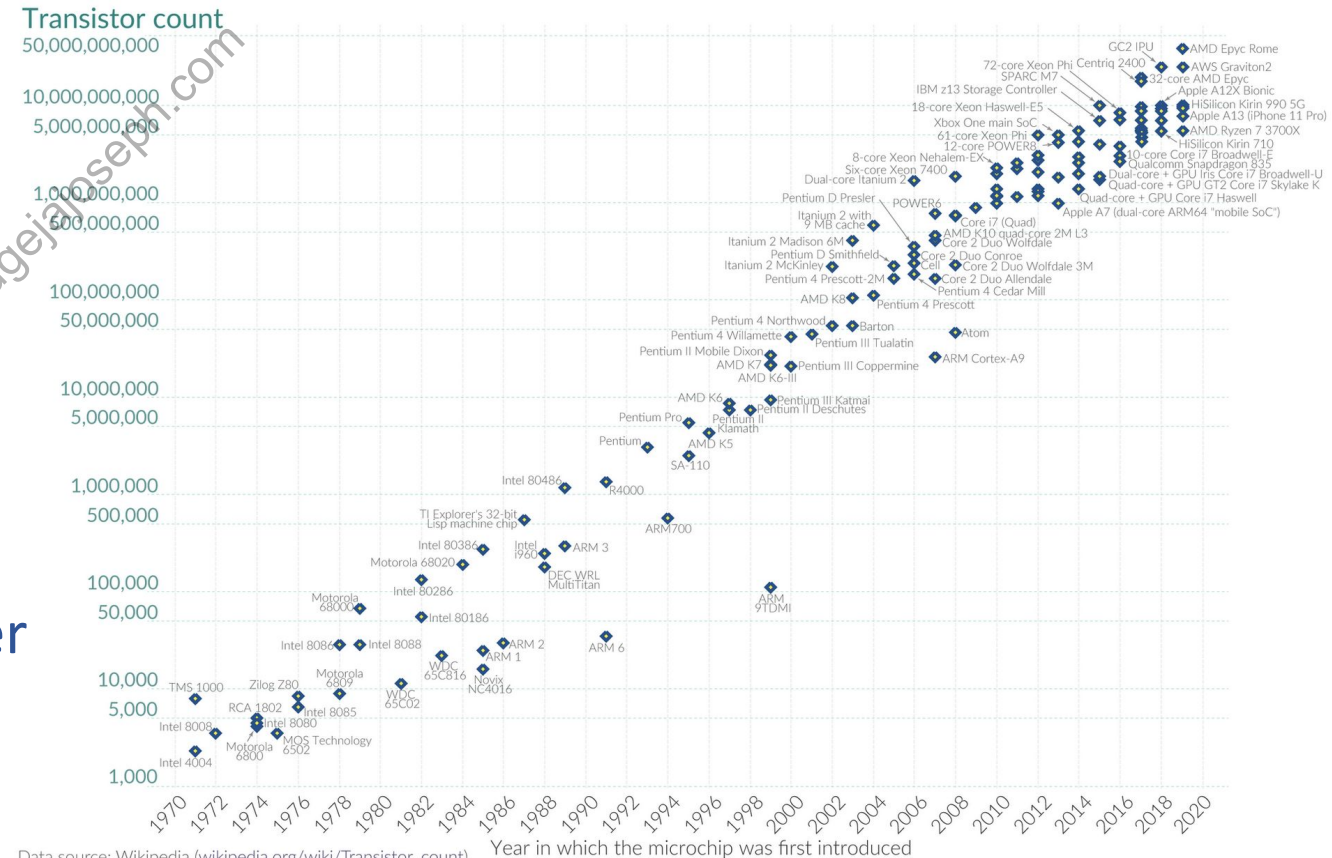


What is Cryptanalysis?

- *Cryptanalysis* is the practice of uncovering flaws within cryptosystems.
- E.g., cryptanalysis has conclusively demonstrated that significant flaws exist in the WEP algorithm.

Brute-Force Attack against Symmetric Ciphers

- Treats the cipher as a black-box
 - Attack: $\forall k_i \in K : d_{k_i}(y) \stackrel{?}{=} x$
 - Requires (at least) 1 plaintext-ciphertext pair (x, y)
 - What key length do we need ?
 - Short-term vs long-term protection
 - Symmetric vs asymmetric crypto
 - Threat and attack model
 - Moore's law
 - etc.
 - A long key space does not help if other attacks are possible.
- 
- The scatter plot shows the exponential growth of transistor counts in microprocessors over time. The y-axis is a logarithmic scale for 'Transistor count' ranging from 5,000 to 50,000,000,000. The x-axis represents time from 1970 to 2010. Data points are labeled with processor names and years, showing a consistent upward trend.
- | Processor | Year | Approximate Transistor Count |
|----------------|------|------------------------------|
| Intel 8005 | 1972 | 600 |
| TMS 1000 | 1974 | 6,000 |
| RCA 1802 | 1975 | 6,000 |
| Intel 8080 | 1976 | 6,000 |
| Zilog Z80 | 1979 | 60,000 |
| Intel 8085 | 1982 | 60,000 |
| Intel 8088 | 1982 | 290,000 |
| Intel 8086 | 1982 | 290,000 |
| Motorola 6800 | 1976 | 6,000 |
| Motorola 6801 | 1978 | 6,000 |
| Motorola 6802 | 1980 | 6,000 |
| Motorola 6805 | 1982 | 6,000 |
| Motorola 6808 | 1983 | 6,000 |
| Motorola 6809 | 1984 | 6,000 |
| Motorola 68000 | 1985 | 12,000,000 |
| Motorola 68010 | 1986 | 12,000,000 |
| Motorola 68020 | 1987 | 12,000,000 |
| Motorola 68030 | 1988 | 12,000,000 |
| Motorola 68040 | 1989 | 12,000,000 |
| Motorola 68060 | 1990 | 12,000,000 |
| Motorola 68080 | 1991 | 12,000,000 |
| Motorola 68090 | 1992 | 12,000,000 |
| Motorola 68010 | 1993 | 12,000,000 |
| Motorola 68020 | 1994 | 12,000,000 |
| Motorola 68030 | 1995 | 12,000,000 |
| Motorola 68040 | 1996 | 12,000,000 |
| Motorola 68050 | 1997 | 12,000,000 |
| Motorola 68060 | 1998 | 12,000,000 |
| Motorola 68070 | 1999 | 12,000,000 |
| Motorola 68080 | 2000 | 12,000,000 |
| Motorola 68090 | 2001 | 12,000,000 |
| Motorola 68010 | 2002 | 12,000,000 |
| Motorola 68020 | 2003 | 12,000,000 |
| Motorola 68030 | 2004 | 12,000,000 |
| Motorola 68040 | 2005 | 12,000,000 |
| Motorola 68050 | 2006 | 12,000,000 |
| Motorola 68060 | 2007 | 12,000,000 |
| Motorola 68070 | 2008 | 12,000,000 |
| Motorola 68080 | 2009 | 12,000,000 |
| Motorola 68090 | 2010 | 12,000,000 |
| Motorola 68010 | 2011 | 12,000,000 |
| Motorola 68020 | 2012 | 12,000,000 |
| Motorola 68030 | 2013 | 12,000,000 |
| Motorola 68040 | 2014 | 12,000,000 |
| Motorola 68050 | 2015 | 12,000,000 |
| Motorola 68060 | 2016 | 12,000,000 |
| Motorola 68070 | 2017 | 12,000,000 |
| Motorola 68080 | 2018 | 12,000,000 |
| Motorola 68090 | 2019 | 12,000,000 |
| Motorola 68010 | 2020 | 12,000,000 |
| Motorola 68020 | 2021 | 12,000,000 |
| Motorola 68030 | 2022 | 12,000,000 |
| Motorola 68040 | 2023 | 12,000,000 |
| Motorola 68050 | 2024 | 12,000,000 |
| Motorola 68060 | 2025 | 12,000,000 |
| Motorola 68070 | 2026 | 12,000,000 |
| Motorola 68080 | 2027 | 12,000,000 |
| Motorola 68090 | 2028 | 12,000,000 |
| Motorola 68010 | 2029 | 12,000,000 |
| Motorola 68020 | 2030 | 12,000,000 |



Data source: Wikipedia ([wikipedia.org/wiki/Transistor_count](https://www.wikipedia.org/wiki/Transistor_count))
OurWorldinData.org – Research and data to make progress against the world's largest problems. Licensed under CC-BY by the authors Hannah Ritchie and Max Roser.

Black-Box Attack Models

- **Ciphertext-only attack:** An attacker observes ciphertexts without having any knowledge of the corresponding plaintext or the method used to select the plaintext.
- **Known-plaintext attack:** An attacker observes ciphertexts and possesses information about the associated plaintext.
- **Chosen-plaintext attack:** An attacker can request encryption for specific plaintexts of their choice and observe the resulting ciphertexts.
- **Chosen-ciphertext attack:** An attacker can both encrypt and decrypt data.



Affine Cipher

- Extension of the shift cipher; rather than just adding the key to the plaintext, we also multiply by the key
- The key consists of two parts: $k = (a, b)$, which has the restriction: $\gcd(a, 26) = 1$

Let $k, x, y \in \{0, 1, \dots, 25\}$

- Encryption: $y = e_k(x) \equiv (a \cdot x + b) \bmod 26$
- Decryption: $x = d_k(y) \equiv a^{-1} \cdot (y - b) \bmod 26$

What is the keyspace of the affine cipher?

- Several attacks are still possible!

Course Overview

Introduction and
Course Overview

Randomness

Stream Ciphers

Block Ciphers
(DES, AES)

Other Aspects of
Block Ciphers

Public-Key
Cryptography

The RSA and
Public-Key
Cryptosystems

Digital Signatures

Hash Functions
and MAC

Key Establishment
and Practical
Applications

Course Summary

Practical Components

- Basic programming code will be utilized during most of the lectures.
- Online tools and resources will be employed to identify and address design or implementation weaknesses.
- Theoretical concepts will be applied to real-world industry scenarios, whenever applicable.